#### Logic Programming, 16-17

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- General problem: given an initial state  $S_i$  and a final state  $S_f$ , find a path between these two states making sure each transition corresponds to a valid move.
- Example: Hanoi
  - blocks need to be moved one at a time.
  - ▶ a block can only be moved if there is nothing on its top.
- To find the solution (sequence of moves), we need to be able to transform our initial blocks state in the final blocks state.

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#### Hanoi example

• Exemplo:



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- what path to follow to transform config 1 in config 2?
- how to find this path?

### Hanoi example

- We will explore alternative paths till finding the solution.
- For example, after placing C in the floor, we have the alternatives:
  - ▶ place A in the floor, **OR**
  - ▶ place A on top of C, **OR**
  - ▶ place C on top of A.
- Two types of concepts:
  - ▶ situations (states, nodes, configurations).
  - possible movements (actions, operators, transformations) which can transform one state in another state.

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#### Search Problems

- Summarizing:
  - ▶ state space.
  - initial state.
  - ▶ final state (goal).
  - operator that can transform one state in the next.
- Optimization problem: find the path with minimum cost.

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- Representation of a state space in Prolog: s(X,Y) ou s(X,Y,C), with C a cost of transitioning from state X to Y.
- s(X,Y) is true is there is a legal/possible movement from X to Y.
- In the blocks problem (Hanoi), a state can be represented by a list of stacks. Each stack, byt its turn can be represented by a list whose first element is the block on the top of the stack.

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# Search Problems: defining a transition from X to Y example

- Initial state: [[c,a,b],[],[]]
- Final state: any set of stacks that contains one of the stacks with the blocks ordered:
  - [[a,b,c],[],[]]
  - [[],[a,b,c],[]]
  - [[],[],[a,b,c]]
- Given a state, to find the next state, we use the following rule: St2 is the next state after St1, if there are two stacks Stk1 and Stk2, with the block on top of Stk1 being moved to Stk2.

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# Search Problems: defining a transition from X to Y example

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% mv Top1 to Stk2 em St2 s(Stacks,[Stk1,[Top1|Stk2]|Otherstacks]) :-% [Top1|Stk1] is a stack in St1 del([Top1|Stk1],Stacks,Stacks1), % Stk2 is a stack in St1 del(Stk2,Stacks1,Otherstacks).

## Search Problems: Hanoi

- goal (final state):
   goal(Estado) :- member([a,b,c],Estado).
- search predicate (can be implemented using any search predicate: dfs, bfs etc:

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solve(Initial,Final).

• Query: ?- solve([[c,a,b],[],[]],Solution).

• Depth-first search (dfs):
solve(N,[N]) :- goal(N).
solve(N,[N|Sol1]) :s(N,N1), % the implementation of s/2
solve(N1,Sol1).% depends on the problem

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• Note: this program does not prevent cycles.

• Iterative Deepening: (needs extra argument: depth limit)
solve(N,[N],\_) :- goal(N).
solve(N,[N|Sol1],ProfMax) :ProfMax > 0,
s(N,N1),
NewMax is ProfMax - 1,
solve(N1,Sol1,NewMax).

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• Breadth-first Search (BFS):

bfs(Initial,Final) :- solve([[Initial]],Final). solve([[N|Path]|\_],[N|Path]) :- goal(N). solve([Path|Paths],Solution) :extend(Path.NewPaths). conc(Paths, NewPaths, Paths1), solve(Paths1,Solution). extend([Node|Path],NewPaths) :bagof([NewNode,Node|Path], (s(Node, NewNode), \+ member(NewNode, [Node | Path])), NewPaths), !. extend(Path,\_). % node has no successor.

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