# Tree Search for the Recursive Circle Packing Problem 

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# Outline 

Recursive Circle Packing Origin of the RCPP Informal description MINLP formulation

Algorithms
Greedy construction
Generating positions
Semi-Greedy construction
Local Search
Depth-First Tree Search Monte-Carlo Tree Search

Computational Results
Conclusion

## Next up...

## Recursive Circle Packing

## Algorithms

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## Origin of the problem: industrial setting

- A company produces hollow tubes of various radii
- Orders are sent to customers in containers
- All tubes have the length of a container
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- Currently: solution is constructed...MANUALLY (what?!? O_o)
- Very tedious and error prone
- Production engineers' time is expensive
- We can surely do better


## Problem description (simplified)

Given:

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- external radius
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## Solution:

- a list of packed tubes and their positions


## A mathematical model: parameters and variables

Parameters for describing an instance

- width $W$ and height $H$ of the rectangular container;
- set of tubes $\mathcal{N}$
- for each tube $i \in \mathcal{N}$ :
- external radius $r_{i}^{\text {ext }}$
- internal radius $r_{i}^{\text {int }}$
- value $v_{i}$


## Variables

- $\left(x_{i}, y_{i}\right) \in \mathbb{R}^{2}$, position of the center of each tube $i$
- $p_{i} \in\{0,1\}, p_{i}=1$ if tube $i$ is placed directly in the container
- $q_{i j} \in\{0,1\}, q_{i j}=1$ if tube $i$ is placed directly inside tube $j$


## A mathematical model: objective function

$$
\operatorname{maximize} V=\sum_{i \in \mathcal{N}} v_{i} \times\left(p_{i}+\sum_{j \in \mathcal{N}} q_{i j}\right)
$$

# A mathematical model: constraints 

Tubes cannot be in multiple places

$$
p_{i}+\sum_{j} q_{i j} \leq 1, \quad \forall i \in \mathcal{N}
$$

Tubes must stay inside the container

$$
\begin{array}{ll}
r_{i}^{\text {ext }} \leq x_{i} \leq W-r_{i}^{\text {ext }}, \quad \forall i \in \mathcal{N} \\
r_{i}^{\text {ext }} \leq y_{i} \leq H-r_{i}^{\text {ext }}, \quad \forall i \in \mathcal{N}
\end{array}
$$

A mathematical model: more constraints
Tubes $i$ and $j$ inside the container cannot overlap

$$
\left\|x y_{i}-x y_{j}\right\|_{2} \geq r_{i}^{\text {ext }}+r_{j}^{\text {ext }}-M \times\left(2-p_{i}-p_{j}\right), \quad \forall i, j \in \mathcal{N}
$$

Same goes for tubes $i$ and $j$ inside the same larger tube $k$

$$
\left\|x y_{i}-x y_{j}\right\|_{2} \geq r_{i}^{\text {ext }}+r_{j}^{\text {ext }}-M \times\left(2-q_{i k}-q_{j k}\right), \quad \forall i, j, k \in \mathcal{N}
$$

Tube $i$ inside $j$ must stay within $j$

$$
\left\|x y_{i}-x y_{j}\right\|_{2} \leq r_{j}^{\text {int }}-r_{i}^{\text {ext }}+M \times\left(1-q_{i j}\right), \quad \forall i, j \in \mathcal{N}
$$

( $M$ disables constraints when any variable inside parenthesis is zero)

## A mathematical model: limitations

- Previous model is exact
- Non-linear
- Very hard to solve
- Quadratic number of variables
- Cubic number of constraints
- Let's take a look at some practical solutions


# Next up... 

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## A possible greedy construction

```
def greedy solution(C, T):
    O = {C}
    while O != {}:
        o = argmin(O, key=free_space)
        repeat:
                if o has no positions:
                O.remove(o)
                break
            t = argmax(T, key=(erad, value, irad))
            P = o.positions_for(t)
            p = argmin(P, key=(y_coord, x_coord))
            o.insert(t, p)
            O.add(t)
            T.remove(t)
            o = t
return C
```


## Generating positions

What is the set of positions for a tube?
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In reality, this is an usually an infinite set.
Reduce this to a finite set of possible positions

- the corners of the container
- positions touching another tube and a wall of the container
- positions touching (at least) two other tubes
- the bottom of the outer tube (telescoping; initial tube)

Generating positions


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- $\mathrm{SG}=$ greedy construction + probabilistic choice + repetition
- Probabilistic choice: position of tube
- Optimization by repetition with different RNG seeds
- Could hardly be simpler


## Local Search

- Not easy to define a proper (finite) neighborhood
- Moving one tube will likely cause overlaps
- May not be trivial to restore feasibility


## Depth-First Tree Search

- Complete enumeration of the search space (if enough time is allowed)
- Avoids repeated solutions
- Very fast
- Low memory requirements
- First decisions are (in most cases) never changed
- Performance extremely dependent on a branch ordering heuristic


## Monte-Carlo Tree Search

## Some facts

- Tree search algorithm mostly employed in game playing
- Asymmetrically constructs a game tree
- Focuses on most promising* branches
- Uses Monte-Carlo simulations to estimate value of nodes
- Each node maintains basic statistics (\# of sims and \# of wins)
- Requires little/no domain-specific knowledge, but benefits from it
- An iteration consists of...


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Selection starting from the root, pick a node to expand Expansion create one (or more) children of the selected node Simulation make a simulation from each new node Backpropagation using the results of the simulations, update the statistics on each node in the path up to the root

## Monte-Carlo Tree Search



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Backpropagation


# UCT: Upper Confidence Bound 1 applied to Trees 

$$
U C T_{n}=X_{n}+c \cdot \sqrt{\frac{\ln N_{P_{n}}}{N_{n}}}
$$

- Formula for selecting the "best" child (selection step)
- Most popular variant of MCTS
- UCT formula consists of two components:

Exploitation prefers nodes with best known values
Exploration prefers nodes that have few simulations

- $X_{n}$ is assumed to be in $[0,1]$


## Adapting UCT for optimization

Normalization

$$
X_{n}=\frac{e^{1-\frac{z_{n}^{*}-z^{*}}{w^{*}}-2^{*}}-1}{e-1}
$$

- $X_{n} \in[0,1] \checkmark$ UCT-approved
- uses both $z^{*}$ and $w^{*}$ to assess how good a value is
- avoids scale issues with objective function values


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Weighting exploration with $\overline{X_{n}}$

$$
E_{n}=\overline{X_{n}} \cdot c \cdot \sqrt{\frac{\ln N_{P_{n}}}{N_{n}}}
$$

- use mean to help guide the search
- assign less time to branches with worse mean score


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## Experiment

- 6 instances
- 3,5 , and 16 different types of tubes
- 2 container sizes: large and small (= large/2)
- Competing algorithms

> DFS Depth-First Tree Search
> SG repeated Semi-Greedy construction MCTS Monte Carlo Tree Search

- Software implemented in Python 2.X, run in PyPy

Results: large instances

|  |  | large03 | large05 | large16 |
| :---: | :---: | :---: | :---: | :---: |
| DFS* $^{*}$ |  | 3570033 | 3720124 | 18492283 |
|  | $\min$ | 3660028 | 3810114 | 22381890 |
| SG** $^{* *}$ | avg | 3660029.3 | 3822105.7 | 22548874.4 |
|  | $\max$ | 3660030 | 3840093 | 22851844 |
|  | $\min$ | 3660029 | 4050053 | 24241737 |
| MCTS** $^{* *}$ | $\operatorname{avg}$ | $\mathbf{3 6 6 0 0 3 1 . 3}$ | $\mathbf{4 0 9 8 0 5 2 . 3}$ | $\mathbf{2 4 8 4 2 6 8 5 . 8}$ |
|  | $\max$ | 3660034 | 4140048 | 25451624 |

* result of 1 run of 600 s
** results of 10 independent runs of 600s


## Results: small instances

|  |  | small03 | small05 | small16 |
| :---: | :---: | :---: | :---: | :---: |
| DFS* $^{*}$ |  | 900000 | 1090000 | 9540056 |
| SG** $^{* *}$ | $\min$ | 940000 | 1090000 | 9790035 |
|  | avg | 940000.0 | 1090000.0 | 9820032.6 |
|  | $\max$ | 940000 | 1090000 | 9840031 |
|  |  | Min | 940000 | 1120000 |
| MCTS** $^{* *}$ | avg | $\mathbf{9 5 6 0 0 0 . 1}$ | $\mathbf{1 1 2 0 0 0 0 . 0}$ | $\mathbf{1 0 6 4 3 0 0 3 4}$ |
|  | $\max$ | 960000 | 1120000 | 10700041 |

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## Contributions

- Definition of the RCPP
- Non-linear formulation: not usable in practice
- Adaptation of MCTS/UCT
- make $X_{n}$ independent of problem scale
- use mean performance to guide exploration
- Interesting results


## Conclusion

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Future work

- Compare with more challenging opponents
- Apply MCTS to other problems, e.g., MIP
- Add some mechanism to discard nodes when the tree grows too large (beam search)


# Thank you! 



