Hybrid Computing



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Why Hybrid?

- Eliminates domain decomposition at node
- Automatic coherency at node
- Lower memory latency and data movement within node
- Can synchronize on memory instead of barrier





Why Hybrid? (cont 1)

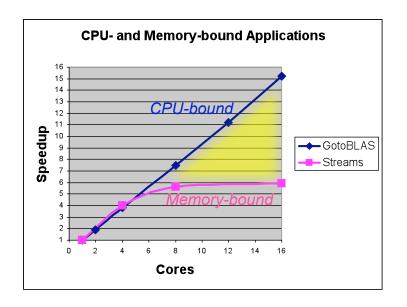
 Only profitable if on-node aggregation of MPI parallel components is faster as a single SMP algorithm (or a single SMP algorithm on each socket).





Hybrid - Motivation

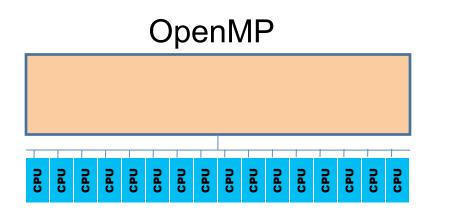
- Load Balancing
- Reduce Memory Traffic

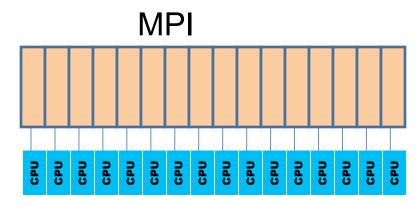




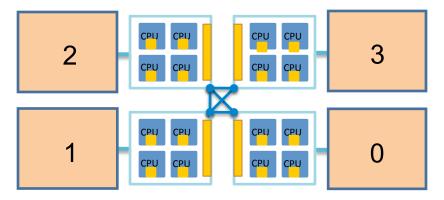


Node Views





Process-Affinity Memory-Allocation







NUMA Operations

- Where do threads/processes and memory allocations go?
- If Memory were completely uniform there would be no need to worry about these two concerns. Only for NUMA (non-uniform memory access) is (re)placement of processes and allocated memory (NUMA Control) of importance.
- Default Control: Decided by policy when process exec'd or thread forked, and when memory allocated. Directed from within Kernel.

NUMA CONTROL IS MANAGED BY THE KERNEL.

NUMA CONTROL CAN BE CHANGED WITH NUMACLT.





NUMA Operations

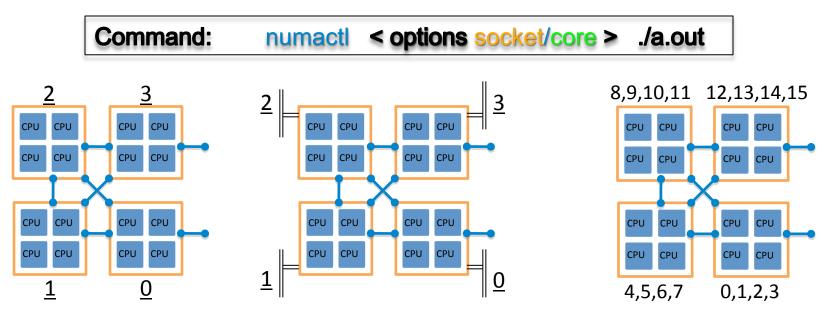
- Ways Process Affinity and Memory Policy can be changed:
 - Dynamically on a running process (knowing process id)
 - At process execution (with wrapper command)
 - Within program through F90/C API
- Users can alter Kernel Policies
 (setting Process Affinity and Memory Policy == PAMPer)
 - Users can PAMPer their own processes.
 - Root can PAMPer any process.
 - Careful, libraries may PAMPer, too!





NUMA Operations

 Process Affinity and Memory Policy can be controlled at <u>socket</u> and <u>core</u> level with <u>numactl</u>.



Process: Socket References
process assignment
-N

Memory: Socket References

memory allocation

—I —i --peferred —m

(local, interleaved, pref., mandatory)

Process: Core References core assignment

—C.





Modes of MPI/Thread Operation

- SMP Nodes
 - Single MPI task launched per node
 - Parallel Threads share all node memory, e.g 16 threads/ node on Ranger.
- SMP Sockets
 - Single MPI task launched on each socket
 - Parallel Thread set shares socket memory, e.g. 4 threads/socket on Ranger
- MPI Cores
 - Each core on a node is assigned an MPI task.
 - (not really hybrid, but in master/slave paradigm master could use threads)

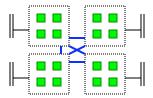




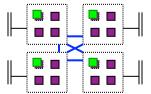
Modes of MPI/Thread Operation

Pure MPI Node Pure SMP Node

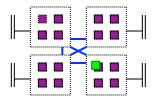
16 MPI Tasks



4 MPI Tasks 4Threads/Task



1 MPI Tasks 16 Threads/Task



Master Thread of MPI Task

MPI Task on Core

Master Thread of MPI Task

Slave Thread of MPI Task





SMP Nodes

Hybrid Batch Script 16 threads/node

- Make sure 1 task is created on each node
- Set total number of cores (nodes x 16)
- Set number of threads for each node
- PAMPering at job level
 - Controls behavior for ALL tasks
 - No simple/standard way to control thread-core affinity

job script (Bourne shell)	job script (C shell)
#! -pe 1way 192	#! -pe 1way 192
export OMP_NUM_THREADS=16	setenv OMP_NUM_THREADS 16
ibrun numactl -i all ./a.out	ibrun numactl –i all ./a.out





SMP Sockets

Hybrid Batch Script 4 tasks/node, 4 threads/task

- Example script setup for a square (6x6 = 36) processor topology.
- Create a task for each socket (4 tasks per node).
- Set total number of cores allocated by batch (nodes x 16 cores/node).
- Set actual number of cores used with MY_NSLOTS.
- Set number of threads for each task
- PAMPering at task level
 - Create script to extract rank for numactl options, and a out execution (TACC MPI systems always assign sequential ranks on a node.
 - No simple/standard way to control thread-core affinity

job script (Bourne shell)	job script (C shell)
#! -pe 4way 48	 #! -pe 4way 48
•••	•••
export MY_NSLOTS =36	setenv MY_NSLOTS 36
export OMP_NUM_THREADS=4	setenv OMP_NUM_THREADS 4
ibrun numa.csh	ibrun numa.sh





SMP Sockets

Hybrid Batch Script 4 tasks/node, 4 threads/task

numa.sh

#!/bin/bash

export MV2_USE_AFFINITY=0 export MV2_ENABLE_AFFINITY=0

#TasksPerNode

TPN=`echo \$PE | sed 's/way//'`
[! \$TPN] && echo TPN NOT defined!
[! \$TPN] && exit 1

socket=\$((\$PMI_RANK % \$TPN))

numactl -N \$socket -m \$socket ./a.out

numa.csh

#!/bin/tcsh

setenv MV2_USE_AFFINITY 0 setenv MV2_ENABLE_AFFINITY 0

#TasksPerNode

set TPN = `echo \$PE | sed 's/way//' if(! \${%TPN}) echo TPN NOT defined! if(! \${%TPN}) exit 0

@ socket = \$PMI_RANK % \$TPN

numactl -N \$socket -m \$socket ./a.out





Hybrid – Program Model

- Start with MPI initialization
- Create OMP parallel regions within MPI task (process).
 - Serial regions are the master thread or MPI task.
 - MPI rank is known to all threads
- Call MPI library in serial and parallel regions.
- Finalize MPI



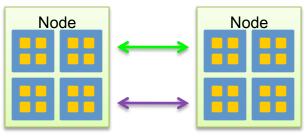


```
MPI_Init
...
MPI_call
OMP Parallel
...
MPI_call
...
end parallel
...
MPI_call
...
MPI_call
```



MPI with OpenMP -- Messaging

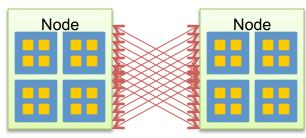
Single-threaded messaging



rank to rank

MPI from serial region or a single thread within parallel region

Multi-threaded messaging



rank-thread ID to any rank-thread ID

MPI from multiple threads within parallel region Requires thread-safe implementation





Threads calling MPI

- Use MPI_Init_thread to select/determine MPI's thread level of support (in lieu of MPI_Init).
 MPI_Init_thread is supported in MPI2
- Thread safety is controlled by "provided" types: single, funneled, serialized and multiple
 - Single means there is no multi-threading.
 - Funneled means only the master thread calls MPI
 - Serialized means multiple threads can call MPI, but only 1 call can be in progress at a time (serialized).
 - Multiple means MPI is thread safe.
 - Monotonic values are assigned to Parameters:

 MPI_THREAD_SINGLE < MPI_THREAD_FUNNELED < MPI_THREAD_SERIALIZED <
 MPI_THREAD_MULTIPLE





MPI2 MPI Init thread

Syntax:

call MPI_Init_thread(irequired, iprovided, ierr) int MPI_Init_thread(int *argc, char ***argv, int required, int *provided) int MPI::Init_thread(int& argc, char**& argv, int required)

Support Levels	Description
MPI_THREAD_SINGLE	Only one thread will execute.
MPI_THREAD_FUNNELED	Process may be multi-threaded, but only main thread will make MPI calls (calls are "funneled" to main thread). Default
MPI_THREAD_SERIALIZE	Process may be multi-threaded, any thread can make MPI calls, but threads cannot execute MPI calls concurrently (MPI calls are "serialized").
MPI_THREAD_MULTIPLE	Multiple threads may call MPI, no restrictions.

If supported, the call will return provided = required. Otherwise, the highest level of support will be provided.





Hybrid Coding

Fortran

```
include 'mpif.h'
program hybsimp

call MPI_Init(ierr)
call MPI_Comm_rank (...,irank,ierr)
call MPI_Comm_size (...,isize,ierr)
! Setup shared mem, comp. & Comm

!$OMP parallel do
do i=1,n
<work>
```

```
call MPI_Finalize(ierr)
```

! compute & communicate

end

enddo

```
#include <mpi.h>
int main(int argc, char **argv){
int rank, size, ierr, i;
ierr= MPI_Init(&argc,&argv[]);
ierr= MPI_Comm_rank (...,&rank);
ierr= MPI_Comm_size (...,&size);
//Setup shared mem, compute & Comm
#pragma omp parallel for
  for(i=0; i<n; i++){
    <work>
// compute & communicate
ierr= MPI_Finalize();
```





MPI Call through Master

- MPI_THREAD_FUNNELED
- Use OMP_BARRIER since there is no implicit barrier in master workshare construct (OMP_MASTER).
- All other threads will be sleeping.





Funneling through Master

```
Fortran
include 'mpif.h'
program hybmas
!$OMP parallel
 !SOMP barrier
 !$OMP master
  call MPI_<Whatever>(...,ierr)
 !$OMP end master
 !$OMP barrier
!$OMP end parallel
end
```

```
#include <mpi.h>
int main(int argc, char **argv){
int rank, size, ierr, i;
#pragma omp parallel
 #pragma omp barrier
 #pragma omp master
  ierr=MPI_<Whatever>(...)
 #pragma omp barrier
```





MPI Call within Single

- MPI_THREAD_SERIALIZED
- Only OMP_BARRIER at beginning, since there is an implicit barrier in SINGLE workshare construct (OMP_SINGLE).
- All other threads will be sleeping.

 (The simplest case is for any thread to execute a single mpi call, e.g. with the "single" omp construct. See next slide.)





Serialize through Single

```
Fortran
include 'mpif.h'
program hybsing
call mpi init thread(MPI THREAD SERIALIZED,
                         iprovided, ierr)
!$OMP parallel
 !$OMP barrier
 !$OMP single
  call MPI_<whatever>(...,ierr)
 !$OMP end single
 !!OMP barrier
!$OMP end parallel
end
```

```
#include <mpi.h>
int main(int argc, char **argv){
int rank, size, ierr, i;
mpi init thread(MPI THREAD SERIALIZED,
                             iprovided)
#pragma omp parallel
 #pragma omp barrier
 #pragma omp single
  ierr=MPI <Whatever>(...)
 //pragma omp barrier
```





Overlapping Communication and Work

- Communicate with one or several cores.
- Work with others during communication.
- Need at least MPI_THREAD_FUNNELED support.
- Can be difficult to manage and load balance!





Overlapping Communication and Work

```
Fortran
include 'mpi.h'
program hybover
!$OMP parallel
 if (ithread .eq. 0) then
   call MPI <whatever>(...,ierr)
 else
   <work>
 endif
!$OMP end parallel
end
```

```
#include <mpi.h>
int main(int argc, char **argv){
int rank, size, ierr, i;
#pragma omp parallel
 if (thread == 0){
   ierr=MPI <Whatever>(...)
 if(thread != 0){
   work
```





Thread-rank Communication

- Can use thread id and rank in communication
- Example illustrates technique in multithread "ping" (send/receive) example.





Thread-rank Communication

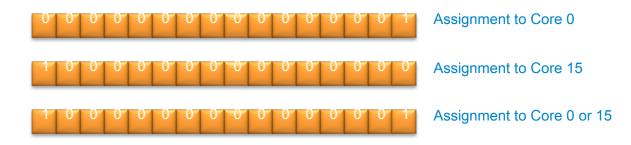
```
call mpi_init_thread( MPI_THREAD_MULTIPLE, iprovided,ierr)
call mpi_comm_rank(MPI_COMM_WORLD, irank, ierr)
call mpi comm size(MPI COMM WORLD, nranks, ierr)
!$OMP parallel private(i, ithread, nthreads)
                                                   Communicate between ranks.
 nthreads=OMP GET NUM THREADS()
 ithread =OMP_GET_THREAD_NUM()
                                                    Threads use tags to differentiate.
 call pwork(ithread, irank, nthreads, nranks...)
 if(irank == 0) then
  call mpi_send(ithread,1,MPI_INTEGER, 1, ithread, MPI_COMM_WORLD, ierr)
 else
  call mpi_recv( j,1,MPI_INTEGER, 0, ithread, MPI_COMM_WORLD, istatus,ierr)
  print*, "Yep, this is ",irank," thread ", ithread," I received from ", j
 endif
!SOMP END PARALLEL
end
```





NUMA in Code

- Scheduling Affinity and Memory Policy can be changed within code through:
 - sched_get/setaffinity
 - get/set_memory_policy
- Scheduling: Bits in a mask are set for assignments.







NUMA in Code

Scheduling Affinity





Conclusion

- Placement and binding of processes, and allocation location of memory are important performance considerations in pure MPI/ OpenMP and Hybrid codes.
- Simple numactl commands and APIs allow users to control process and memory assignments.
- 8-core and 16-core socket systems are on the way, even more effort will be focused on process scheduling and memory location.
- Expect to see more multi-threaded libraries.





References

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