

Bidirectional Data Transformation by Calculation

Hugo Pacheco

HASLab

INESC TEC & University of Minho, Braga, Portugal

Former

55th ToPS Seminar

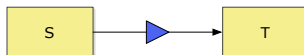
Tokyo - December 18th 2012

Outline

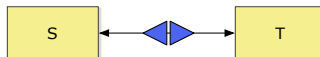
- 1 Introduction
- 2 Efficiency
- 3 Configurability
- 4 Genericity
- 5 Summary

Data Transformations

- data transformations abound in software engineering
- essential to convert data between different formats

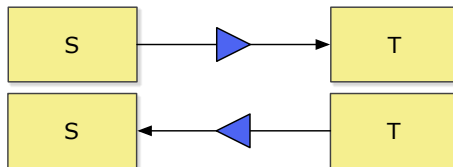


- in real model-driven software engineering scenarios, we often need to run a transformation in both directions



- a **bidirectional transformation (BX)**

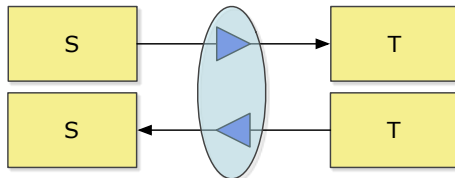
(Ad hoc) Bidirectional Transformations



Manual design: two separate transformations

- expensive
- error-prone
- a maintenance problem

Bidirectional Languages

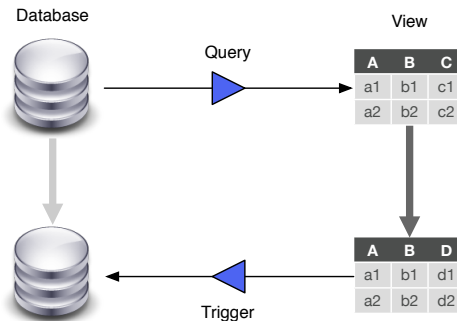


Combinatorial design: the same specification denotes both

- nice syntax
- clean semantics
- compositional

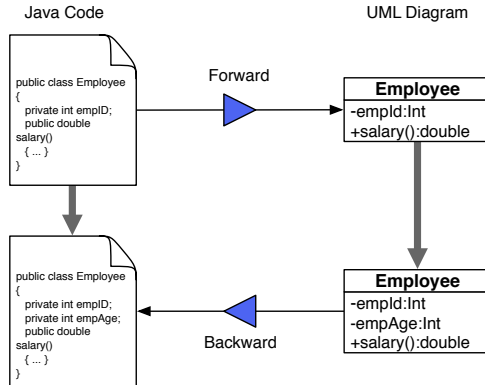
Bidirectional Languages exist for ...

...databases...



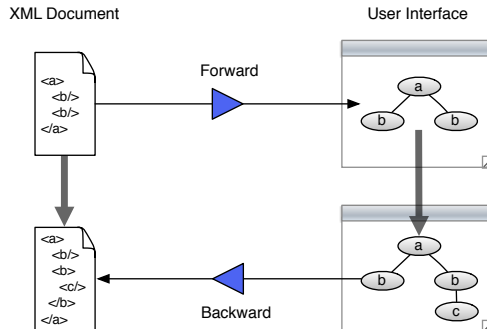
Bidirectional Languages exist for ...

...model-driven software engineering...



Bidirectional Languages exist for ...

...user interfaces...



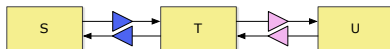
...etc

1

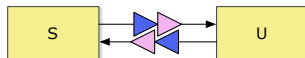
Efficiency

Motivation - Optimization

- combinatorial approaches build complex transformations by composition



- composition \Rightarrow cluttering \Rightarrow inefficiency!



- a serious implementation of BXs needs to be efficient

Question

- how to optimize bidirectional transformations?

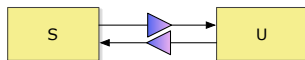


Motivation - Is BX optimization really hard?

- write the two transformations in a language with support for optimization



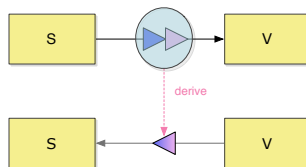
- optimize both independently



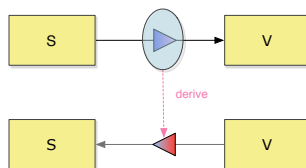
- twice the effort!
- no longer a single program!

Motivation - Is BX optimization really hard?

- write the forward transformation in a language with support for optimization and derive the backward transformation



- optimize the forward transformation and derive the other



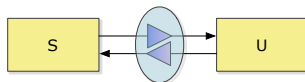
- **different semantics!**

Motivation - A better solution

- write the BX in a language with support for optimization



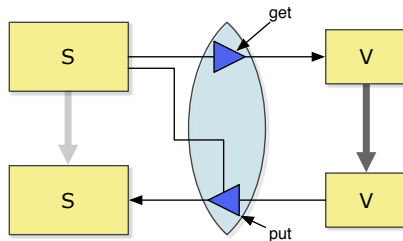
- optimize bidirectional programs



- optimized BX program
- same semantics

Bidirectional Lenses

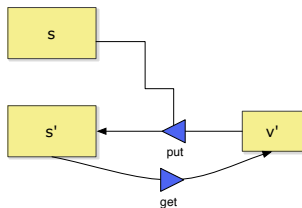
- Foster et al. proposed the framework of lenses [POPL'05, TOPLAS]
- lenses are one of the most popular BX frameworks



The Lens Laws

- PUTGET law

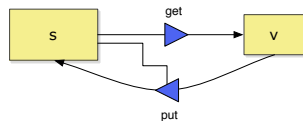
*put must translate
view updates exactly.*



$$get (put\ v'\ s) = v'$$

- GETPUT law

*put must preserve
null view updates.*



$$put (get\ s)\ s = s$$

A Point-free Design

- A data domain (algebraic data types)

data $[a] = [] \mid a : [a]$

data $Tree\ a = Empty \mid Node\ a\ (Tree\ a)\ (Tree\ a)$

- A language syntax

$id : A \rightarrow A$

$\circ : (B \rightarrow C) \rightarrow (A \rightarrow B) \rightarrow (A \rightarrow C)$

$\pi_1 : A \times B \rightarrow A$

$\times : (A \rightarrow C) \rightarrow (B \rightarrow D) \rightarrow (A \times B \rightarrow C \times D)$

- A set of calculation laws

$$f \circ (g \circ h) = (f \circ g) \circ h$$

COMP-ASSOC

$$\pi_1 \circ (f \triangle g) = f \wedge \pi_2 \circ (f \triangle g) = g$$

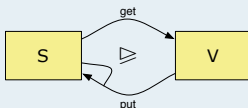
PROD-CANCEL

$$(f \times g) \circ (h \triangle i) = f \circ h \triangle g \circ i$$

PROD-ABSOR

Point-free Lenses

Lens



$$get : S \rightarrow V$$

$$put : V \times S \rightarrow S$$

$$get \circ put = \pi_1$$

$$put \circ (get \triangle id) = id$$

Language of point-free lens combinators

$$Lens ::= id \mid Lens \circ Lens \mid !^c \mid Prod \mid Sum \mid Iso \mid Rec$$

$$Prod ::= \pi_1^b \mid \pi_2^a \mid Lens \times Lens$$

$$Sum ::= Lens \blacktriangledown Lens \mid Lens \nabla_\bullet Lens \mid Lens + Lens \\ \mid inl \nabla Lens \mid Lens \nabla inr$$

$$Iso ::= assocl \mid assocr \mid coassocl \mid coassocr \\ \mid swap \mid coswap \mid distl \mid distr$$

$$Rec ::= in_F \mid out_F \mid F \, Lens \mid (Lens)_F \mid \llbracket Lens \rrbracket_F$$

Point-free Lens Examples

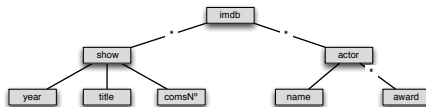
$$\text{length}^a = \llbracket (id + \pi_2^a) \circ out \rrbracket : [A] \triangleright Nat$$

$$\text{map } f = \llbracket in \circ (id + f \times id) \rrbracket : [A] \triangleright [B]$$

$$\text{concat} = \llbracket \dots \rrbracket : [[A]] \triangleright [A]$$


$$ex = \text{map } show \times \text{map } actor$$

$$show = id \times (id \times \text{length} \circ \text{map } (id \times \pi_{Comment}))$$

$$actor = id \times \text{concat} \circ \text{map } \pi_{Awards}$$


Point-free Lens Calculus

- lift the point-free laws to lenses

$$f = g \Leftrightarrow \begin{cases} get_f &= get_g \\ put_f &= put_g \end{cases}$$

- point-free lens laws

$$\pi_1^a \circ (f \times g) = f \circ \pi_1^{createf\ a}$$

PROD-CANCEL

$$(f \nabla g) \circ (h + i) = f \circ h \nabla g \circ i$$

SUM-ABSOR

$$f \circ \langle g \rangle_F = \langle h \rangle_F \Leftarrow f \circ g = h \circ F\ f$$

CATA-FUSION

- fusion examples

$$length^a \circ map\ f = length^{createf\ a}$$

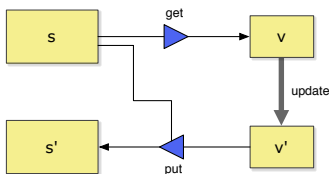
LENGTH-MAP

$$concat \circ map\ f = concatMap\ f$$

CONCAT-MAP

Point-free Lens Calculus - Tupling

- *put* has redundant computations of *get*
- fuse *get* and *put* into a single function - Takeichi [IFP'09]



$$get : S \rightarrow V$$

$$put : S \rightarrow V \rightarrow S$$

$$get \triangle put : S \rightarrow V \times (V \rightarrow S)$$

- Fokkinga's Mutu Tupling theorem

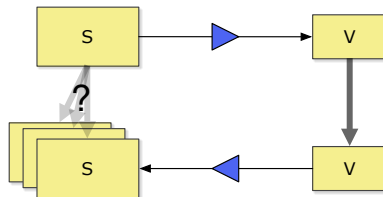
$$f \triangle g = (\phi \triangle \psi)_F \iff \begin{cases} f = \phi \circ F(f \triangle g) \circ out_F \\ g = \psi \circ F(f \triangle g) \circ out_F \end{cases}$$

2

Configurability

Motivation - Configurability

- for non-bijective transformations, an update may have many corresponding updates



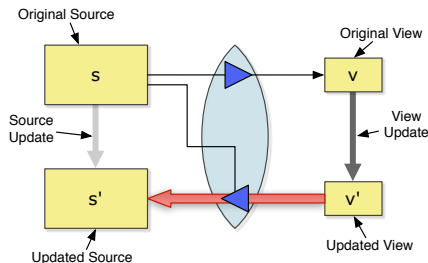
- our point-free lens language provides one possible update
- **may not match the user's intentions!**

Question

- how to allow users to **choose** a suitable update?

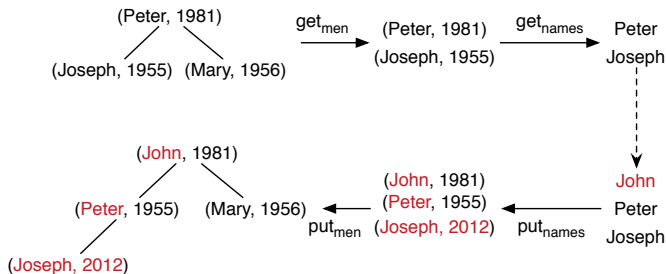
State-based Lenses

State-based framework: *put* takes the modified view state



- no information about the actual update
- *put* has to “guess” the intended change of the update

State-based Lens Example



data $\text{Tree } a = \text{Empty} \mid \text{Node } a \ (\text{Tree } a) \ (\text{Tree } a)$

type $\text{Person} = (\text{Name}, \text{Birth})$

$\text{men} : \text{Tree Person} \triangleright [\text{Person}]$

$\text{men Empty} = []$

$\text{men (Node } p \ f \ m) = p : \text{men } f$

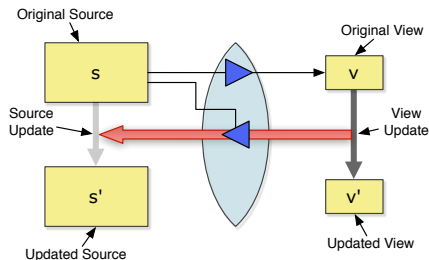
$\text{names} : [\text{Person}] \triangleright [\text{Name}]$

$\text{names} = \text{map } \pi_1^{\text{const } 2012}$

- positional behavior: birth years? Mary?

Operation-based Lenses

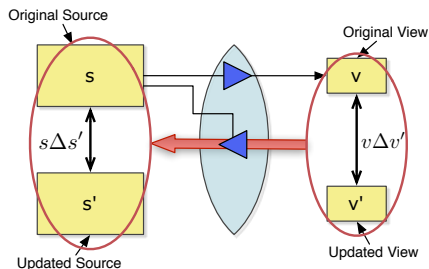
Operation-based framework: *put* takes a representation of the view update



- some knowledge about the actual update
- *put* can infer the intended change

Delta-based Lenses

Delta-based framework: *put* takes a delta



- Diskin et al. proposed an abstract framework of **delta lenses** [JOT]
- deltas model “change”
- delta = description of an update

Point-free Delta Lenses

- we instantiate the abstract framework of Diskin et al.
- we introduce a notion of deltas for algebraic data types

① decompose: shape + data

$l : [\text{Char}]$

x 0

y 1

z 2

② $a\Delta b = \text{partial function}$

$a : [\text{Char}]$

x 0

y 1

z 2

$b : [\text{Char}]$

w 0

x 1

z 2

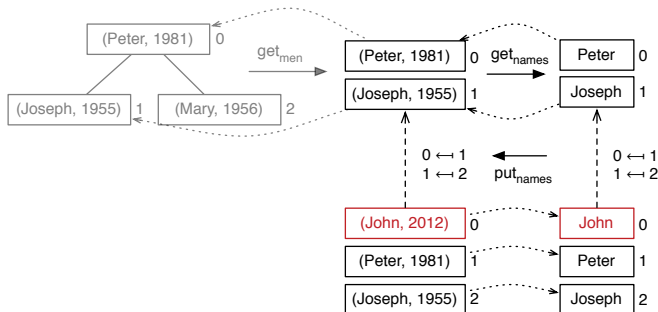


data $[a] = [] \mid a : [a]$

- we define a language of point-free delta lenses
- we tailor our previous lens language to consider deltas

Delta-based Example (Mapping)

- for mapping lenses, that preserve the shape

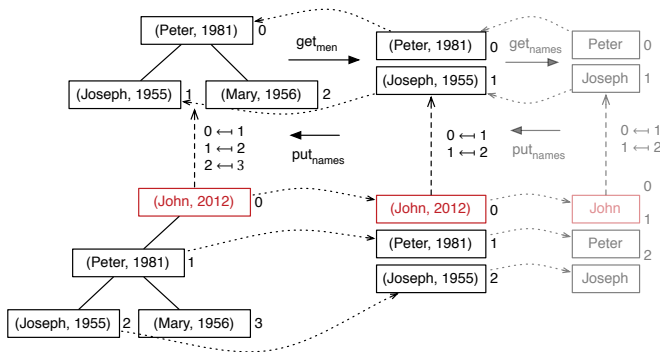


$names : [Person] \triangleright [Name]$

- data alignment

Delta-based Example (Reshaping)

- for reshaping lenses, that preserve the data

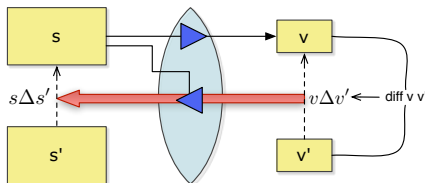


$men : Tree\ Person \triangleright [Person]$

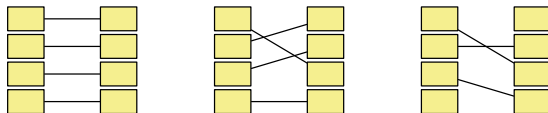
- shape alignment
 - identify shape updates: insertions/deletions
 - propagate shape updates: insertions/deletions

Delta Lens Configurability

- a delta can be calculated from the original and updated states



- user can choose an arbitrary heuristic

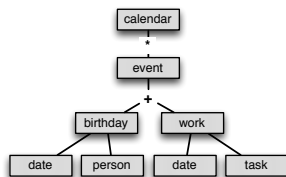


3

Genericity

Motivation - Conciseness

- bidirectional transformations are typically built to match a specific structure, via multiple steps
- collect all event dates in a calendar data format



$calendarDts : Calendar \triangleright Date$

$calendarDts = map\ eventDts$

$eventDts = birthdayDts \nabla \bullet workDts$

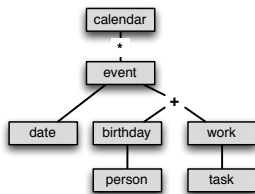
$birthdayDts = \pi_1^{const\ "unknown"}$

$workDts = \pi_1^{const\ ""}$

- **impractical!**
- for a real calendar format (iCal, XML, etc), it would be much more boring

Motivation - Reusability

- for the same high-level transformation, different BXs for different structures (collect all dates)

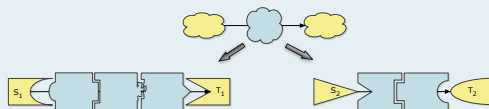


$calendarDts' : Calendar' \triangleright Date$
 $calendarDts' = map\ eventsDts'$
 $eventDts' = \pi_1^{const} (Right\ (Work\ ""))$

- does not support evolution!

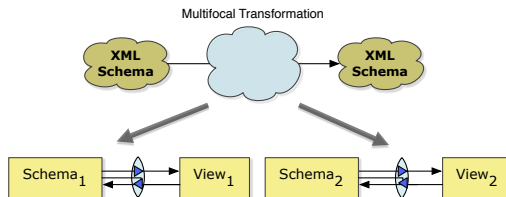
Question

- how to define a transformation in a **concise** and **reusable** way?



The *Multifocal* Language

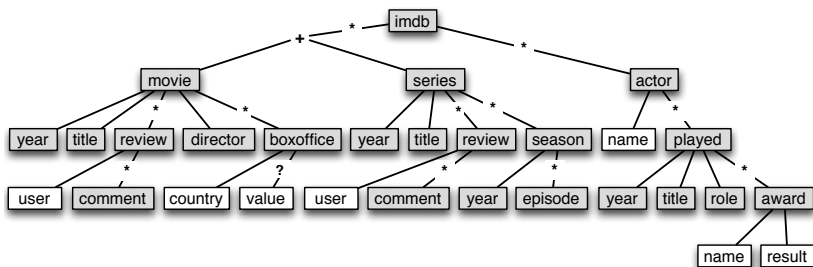
- we propose the *Multifocal* XML transformation language



- Two-level:
 - 1 schema-level transformations as views between XML Schemas
 - 2 model-level transformations as lenses between XML documents
- Strategic: concise specification style (e.g. traversals)
- Bidirectional: underlying document transformations as lenses

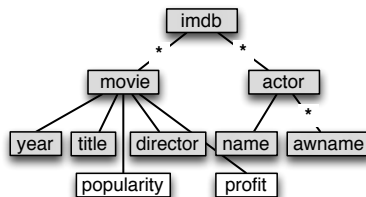
Multifocal Example: XML Views

- source XML Schema modeling a movie database



Multifocal Example: XML Views

- informal XML Schema transformation
 - 1 delete series
 - 2 for each movie:
 - count its popularity (total number of review comments)
 - estimate its profit (sum of the boxoffice values)
 - 3 for each actor, select its name and a list of award names
- view XML Schema



Multifocal Example: XML Views

```
<imdb>
  <movie>
    <year>2003</year>
    <title>Kill Bill: Vol. 1</title>
    <review user="emma">
      <comment>Gorgeous!</comment></review>
    <director>Quentin Tarantino</director>
    <boxoffice country="USA" value="22089322"/>
    <boxoffice country="Japan" value="3521628"/>
  </movie>
  <series><year>2011</year>
    <title>Game of Thrones</title>
    <season><year>2011</year>
      <episode>Winter is Coming</episode>
    </season></series>
  <actor name="Uma Thurman">
    <played><year>2003</year>
      <title>Kill Bill: Vol. 1</title>
      <role>The Bride</role>
      <award name="Saturn" result="Won"/>
    </played></actor>
</imdb>
```

Multifocal Example: XML Views

```
<imdb>
  <movie popularity="1" profit="25610950">
    <year>2003</year>
    <title>Kill Bill: Vol. 1</title>
    <director>Quentin Tarantino</director>
  </movie>
  <actor name="Uma Thurman">
    <awname>Saturn</awname>
  </actor>
</imdb>
```

Multifocal Example: XML Views

```
<imdb>
  <movie> ... </movie>
  <movie popularity="2" profit="15">
    <year>2012</year>
    <title>Sherlock Holmes: Game of Shadows</title>
    <director>Guy Ritchie</director>
  </movie>
  <actor name="Uma Thurman">
    <awname>Saturn Best Actress</awname>
  </actor>
</imdb>
```

Multifocal Example: XML Views

```
<imdb>
  <movie> ... </movie>
  <series> ... </series>
  <movie><year>2012</year>
    <title>Sherlock Holmes: Game of Shadows</title>
    <review user="" comment="" />
    <review user="" comment="" />
    <director>Guy Ritchie</director>
    <boxoffice country="" value="15" />
  </movie>
  <actor name="Uma Thurman">
    <played><year>2003</year>
      <title>Kill Bill: Vol. 1</title>
      <role>The Bride</role>
      <award name="Saturn Best Actress" result="Won"/>
    </played></actor>
</imdb>
```


Multifocal Language: Basic Combinators

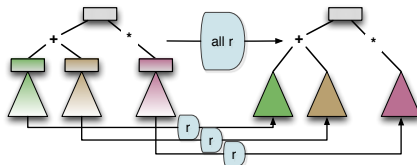
- generic style = concise specification
- strategic combinators

$$\textit{Strat} = \textit{Schema} \rightarrow \textit{Maybe} (\textit{Schema}, \textit{Lens})$$

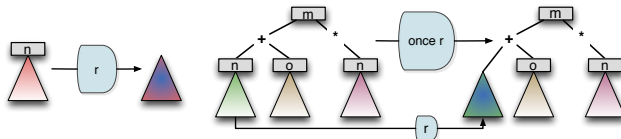
- construct flexible strategies in a compositional way
- basic combinators (in what order? how often?):
 - identity $\text{nop} : \textit{Strat} \rightarrow \textit{Strat}$
 - sequentially $(\gg) : \textit{Strat} \rightarrow \textit{Strat} \rightarrow \textit{Strat}$
 - alternatively $(||) : \textit{Strat} \rightarrow \textit{Strat} \rightarrow \textit{Strat}$
 - repetitively $\text{many} : \textit{Strat} \rightarrow \textit{Strat}$
 - optionally $\text{try} : \textit{Strat} \rightarrow \textit{Strat}$

Multifocal Language: Traversal Combinators

- traversal combinators (at what depth?)
 - apply a strategy to all children



- apply a strategy to all descendants
everywhere : $Strat \rightarrow Strat$
- apply a strategy once at an arbitrary depth

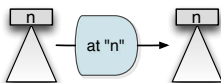


- apply a strategy many times at an arbitrary depth
outermost : $Strat \rightarrow Strat$

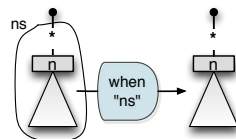
Multifocal Language: Local Combinators

- control the application of certain strategies
- local combinators (under which conditions?)

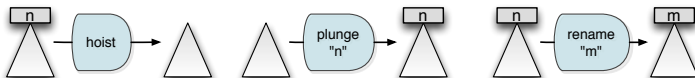
- at a particular element



- at a particular location



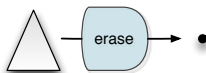
- XML name-based combinators



Multifocal Language: Abstraction Combinators

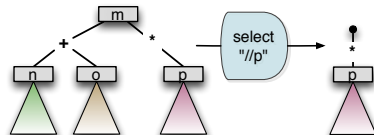
- language for defining XML views
- abstraction combinators (what to delete?)

- erase the current tree (explicit)



- empty tree

- apply an XPath query (implicit)



- 1 specialize the XPath expression ($/m / p$) for the source schema
- 2 convert it to a lens into the query's result type

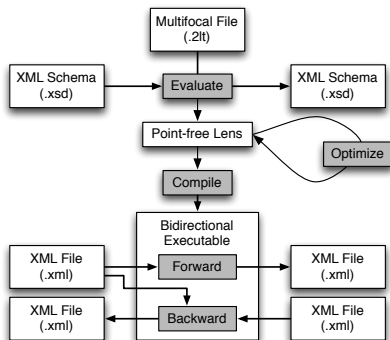
Multifocal Example: Strategic XML Views

- ❶ delete series
- ❷ for each movie:
 - count its popularity (total number of review comments)
 - estimate its profit (sum of the boxoffice values)
- ❸ for each actor, select its name and a list of award names

```
❶ everywhere (try (at "series" erase))  
❷ >> everywhere (try (at "movie" (  
  outermost (when "reviews" (  
    select "count(//comment)" >> plunge "@popularity"))  
  >> outermost (when "boxoffices" (  
    select "sum(//@value)" >> plunge "@profit")))))  
❸ >> everywhere (try (at "actor" (  
  outermost (at "played" (  
    select "award/@name" >> all (rename "awname"))))))
```

The *Multifocal* Framework

- we implement the *Multifocal* framework



- three stages:
 - 1 **evaluate**: XML Schema \Rightarrow XML Schema + lens
 - 2 **optimize (optional)**: lens \Rightarrow optimized lens
 - 3 **compile**: (optimized) lens \Rightarrow executable

Wrapping Up

- ① efficiency
 - point-free lens language
 - algebraic calculus of lenses (fusion, tupling)
- ② configurability
 - point-free delta lens language
 - user-provided alignment heuristics
- ③ genericity
 - *Multifocal* language and framework
 - a *Multifocal* XML schema transformation yields BXs in our lens language that can be optimized
 - reusable: multiple schemas
 - concise: strategic combinators

Libraries and Tools

- implemented in Haskell++
- available online

`http://hackage.haskell.org/package/pointless-lenses`

`cabal install <package>`

`pointless-lenses` (point-free lens library)

`pointless-rewrite` (point-free optimization library)

`pointless-2lt` (strategic two-level lens library)

`multifocal` (*Multifocal* system)

Further Reading



Hugo Pacheco and Alcino Cunha

Generic Point-free Lenses

MPC 2010.



Hugo Pacheco and Alcino Cunha

Calculating with lenses: optimising bidirectional transformations

PEPM 2011.



Hugo Pacheco, Alcino Cunha and Zhenjiang Hu

Delta Lenses over Inductive Types

BX 2012.



Hugo Pacheco and Alcino Cunha

Multifocal: A Strategic Bidirectional Transformation Language for XML Schemas

ICMT 2012.



Alcino Cunha and Hugo Pacheco

Algebraic Specialization of Generic Functions for Recursive Types

Electronic Notes in Theoretical Computer Science, 2011.