Concurrent Programming - Exercises 4 Ismorphism, Trace Equivalence and (Strong) Bisimulation

Consult Chap. 3-1-3,3 e 3.5 of [1].

1. Let

$$P := a!.0 + b!.0$$
$$Q := b!.0 + a!.0$$

determine if the following relations hold

$$\begin{array}{rccc} P & \sim_{tr} & Q \\ \\ a!.P & \sim_{tr} & a!.Q \\ a!.P + a!.P & \sim_{tr} & a!.Q + a!.Q \end{array}$$

and also for \sim_{iso} ?

- 2. Show $(b!.0+c!.0) \sim_{iso} (c!.0+b!.0)$ but $a!.(b!.0+c!0) + a!.(b!.0+c!0) \not\sim_{iso} a!.(b!.0+c!0) + a!.(c!.0+b!.0)$ and conclude that \sim_{iso} is not a congruence.
- 3. Show that for every $P, Q, R \in CCS$

$$P + Q \sim_{tr} Q + P$$

$$(P + Q) + R \sim_{tr} P + (Q + R)$$

$$P + 0 \sim_{tr} P,$$

and $\alpha . (P+Q) \sim_{tr} \alpha . P + \alpha . Q$.

4. Sendo

$$\begin{array}{rcl} CTM & := & coin?.(coffee!.CTM + tea!.CTM) \\ CTM' & := & coin?.coffee!.CTM' + coin?.tea!.CTM' \end{array}$$

Mostra que $CTM \sim_{tr} CTM'$.

- 5. Solve bisimulation problems in http://tinyurl.com/pseuco
- Solve (strong) bisimulation problems in PseuCo Book https://book.pseuco.com/#/read/ equality/workout
- 7. Show that
 - (a) \sim is an equivalence relation.
 - (b) \sim is the largest bisimulation.
 - (c) $s \sim t$ iff for each $\alpha \in Act$

- If $s \xrightarrow{\alpha} s'$ then t' such that $t \xrightarrow{\alpha} t' e s' \sim t'$
- If $t \xrightarrow{\alpha} t'$ then s' such that $s \xrightarrow{\alpha} s'$ and $s' \sim t'$.
- 8. Let $TS = (S, Act, \rightarrow)$ such that $S = \{s_i \mid i \ge 1\} \cup \{t\}, Act = \{a\} \in a = \{(s_i, s_{i+1} \mid i \ge 1\} \cup \{(t, t)\}.$

Show that $s_1 \sim t$ proving that $R = \{(s_i, t) \mid i \geq 1\}$ is a bisimulation.

9. Let P and Q be defined by

$$P := a.P_1 + b.P_2$$
$$P_1 := c.P$$
$$P_2 := c.P$$

and

$$Q := a.Q_1 + b.Q_2$$

$$Q_1 := c.Q_3$$

$$Q_2 := c.Q_3$$

$$Q_3 := a.Q_1 + b.Q_2$$

Show that $P \sim Q$ presenting a bisimulation that contains (P, Q). Draw their LTSs and test in pseuCo.com.

10. Let P and Q be defined by

$$P := a.P_1$$
$$P_1 := b.P + c.P$$

and

$$Q := a.Q_1$$

$$Q_1 := b.Q_2 + c.Q$$

$$Q_2 := a.Q_3$$

$$Q_3 := b.Q + c.Q_2$$

Show that $P \sim Q$ presenting a bisimulation that contains (P, Q).

11. Show that if $P \sim Q$ and $\alpha \in Act, R \in CCS$ and $H \subseteq Com$, then

$$\begin{array}{rcl} \alpha.P & \sim & \alpha.Q \\ P+R & \sim & Q+R \\ R+P & \sim & R+Q \\ P|R & \sim & Q|R \\ R|P & \sim & R|Q \\ P\backslash H & \sim & Q\backslash H \end{array}$$

12. Let

$$P := a.(b.0 + c.0)$$

 $Q := a.b.0 + a.c.0$

Show that $P\;$ and Q are not strongly bisimilar.

13. Show that

$$\begin{array}{rcl} P+Q & \sim & Q+P \\ P+0 & \sim & P \\ (P+Q)+R & \sim & P+(Q+R) \end{array}$$

14. Show that for all $P, Q, R \in CCS$,

$$\begin{array}{rcl} P|Q & \sim & Q|P \\ P|0 & \sim & P \\ (P|Q)|R & \sim & P|(Q|R) \end{array}$$

Show that these are bisimulations

$$\begin{split} &\{(P|Q,Q|P) \mid P,Q \in CCS\}, \\ &\{(P|0,P) \mid P \in CCS\}, \\ &\{((P|Q)|R,P|(Q|R)) \mid P,Q,R \in CCS\}. \end{split}$$

15. Let:

$$\begin{array}{lll} C_0 &:=& inc.C_1\\ C_n &:=& inc.C_{n+1}+dec.C_{n-1}, \mbox{ para }n\geq 1 \end{array}$$

and C := inc.(C|dec.0). Show that the following relation is a bisimulation.

$$\mathcal{R} = \{ (C | \prod_{i=1}^{k} P_i, C_n) | k \ge 0 \land (P_i = 0 \lor P_i = dec.0) \land \text{ the number of } is \text{ with } P_i = dec.0 \text{ is } n \}$$

Consider $(C|\prod_{i=1}^k P_i, C_n) \in \mathcal{R}$. Show that

1. if $C | \prod_{i=1}^{k} P_i \xrightarrow{\alpha} P$ exists Q such that $C_n \xrightarrow{\alpha} Q$ and $(P,Q) \in \mathcal{R}$. 2. if $C_n \xrightarrow{\alpha} Q$ exists P such that $C | \prod_{i=1}^{k} P_i \xrightarrow{\alpha} P$ and $(P,Q) \in \mathcal{R}$. 16. Consider a 1-Buffer.

$$B := put?.get?.B$$

For $n \ge 1$ we can consider a Buffer with capacity n, where B_i^n is a buffer with capacity n with $0 \le i \le n$ elements.

- (a) Verify that $B \sim_{iso} B_0^1$ (draw their LTSs).
- (b) Verify that $B_0^2 \sim B_0^1 | B_0^1$ (draw their LTSs).
- (c) Show that for $n \ge 1$, $B_0^n \sim \underbrace{B_0^1 | B_0^1 | \dots | B_0^1}_{1 + 1 + 1}$.

showing that this is a bisimulation

$$\mathcal{R} = \{ (B_i^n, B_{i_1}^1 | B_{i_2}^1 | \dots | B_{i_n}^1) \mid i_j \in \{0, 1\} \land \sum_{j=1}^n i_j = i \}$$

References

[1] Luca Aceto, Anna Ingólfsdóttir, and Kim Larsen. *Reactive Systems: Modeling, Specification and Verification.* CUP, 2007.