5
The typed λ-calculus

So far we have discussed the type-free λ-calculus. In this chapter we will discuss type systems for the λ-calculus. The initial motivation to define typed versions of the λ-calculus was to avoid paradoxical uses of the untyped calculus [Church, 1940].

5.1 Simple Types

The Curry Type System was first studied in [Curry, 1934] for the theory of combinators. In [Curry and Feys, 1958] this system was modified for the λ-calculus. The definitions and proofs of results in this section can be found in [Barendregt, 1992].

We start by defining the set of types for this system.

Definition 5.1

Let $\mathbb{V}$ be an infinite set of type variables. The *set of simple types*, $\mathbb{T}_C$ is inductively defined from $\mathbb{V}$ in the following way:

$$\begin{array}{rcl} \alpha \in \mathbb{V} & \Rightarrow & \alpha \in \mathbb{T}_C \\ \tau, \tau' \in \mathbb{T}_C & \Rightarrow & (\tau \to \tau') \in \mathbb{T}_C \end{array}$$

Notation. If $\tau_1, \ldots, \tau_n \in \mathbb{T}_C$, then

$$\tau_1 \to \tau_2 \to \cdots \to \tau_n$$

represents

$$(\tau_1 \to (\tau_2 \to \cdots \to (\tau_{n-1} \to \tau_n) \ldots))$$

that is, the type constructor $\to$ is right associative.

Definition 5.2

If x is a term variable in $\mathcal{V}$ and τ is a type in $\mathbb{T}_C$ then:

- A *statement* is of the form $M : \tau$, where the type τ is called the *predicate*, and the variable x is called the *subject* of the statement.
- A *declaration* is a statement where the subject is a term variable.
- A *basis* Γ is a set of declarations where all the subjects are distinct.

A basis where the subjects are pairwise distinct, is called *monovalent*.

Definition 5.3

In the Curry type system, we say that M has type τ given the basis Γ, and write

$$\Gamma \vdash_{\mathcal{C}} M : \tau,$$

if $\Gamma \vdash_{\mathcal{C}} M : \tau$ can be obtained from the following derivation rules:

$$\Gamma \cup \{x : \tau\} \vdash_{\mathcal{C}} x : \tau \qquad (\mathsf{Axiom})$$

$$\frac{\Gamma \cup \{x : \tau_1\} \vdash_{\mathcal{C}} M : \tau_2}{\Gamma \vdash_{\mathcal{C}} \lambda x.M : \tau_1 \to \tau_2} \qquad (\to \mathsf{Intro})$$

$$\frac{\Gamma \vdash_{\mathcal{C}} M : \tau_1 \to \tau_2 \qquad \Gamma \vdash_{\mathcal{C}} N : \tau_1}{\Gamma \vdash_{\mathcal{C}} MN : \tau_2} \qquad (\to \mathsf{Elim})$$

Example 5.4

For the λ-term $(\lambda xy.x)(\lambda x.x)$ the following derivation is obtained in the Curry Simple Type System:

$$\dfrac{\dfrac{\dfrac{\{x : \alpha \to \alpha, y : \beta\} \vdash_{\mathcal{C}} x : \alpha \to \alpha}{\{x : \alpha \to \alpha\} \vdash_{\mathcal{C}} \lambda y.x : \beta \to \alpha \to \alpha}}{\vdash_{\mathcal{C}} \lambda xy.x : (\alpha \to \alpha) \to \beta \to \alpha \to \alpha} \qquad \dfrac{\{x : \alpha\} \vdash_{\mathcal{C}} x : \alpha}{\vdash_{\mathcal{C}} \lambda x.x : \alpha \to \alpha}}{\vdash_{\mathcal{C}} (\lambda xy.x)(\lambda x.x) : \beta \to \alpha \to \alpha}$$

Definition 5.5

If $\Gamma = \{x_1 : \tau_1, \ldots, x_n : \tau_n\}$ is a basis, then:

- Γ is a partial function, with *domain*, denoted $\mathsf{dom}(\Gamma) = \{x_1, \ldots, x_n\}$, and $\Gamma(x_i) = \tau_i$.
- Let $\mathcal{V}_0$ be a set of variables. Then $\Gamma \restriction \mathcal{V}_0 = \{x : \Gamma(x) \mid x \in \mathcal{V}_0\}$.
- We define Γ_x as $\Gamma \setminus \{x : \tau\}$.

Proposition 5.6 (Basic lemmas)

Let Γ be a basis:

- Let Γ' be a basis such that $\Gamma \subseteq \Gamma'$, then

$$\Gamma \vdash_{\mathcal{C}} M : \tau \Rightarrow \Gamma' \vdash_{\mathcal{C}} M : \tau.$$

- If $\Gamma \vdash_{\mathcal{C}} M : \tau$, then $\mathsf{fv}(M) \subseteq \mathsf{dom}(\Gamma)$.
- If $\Gamma \vdash_{\mathcal{C}} M : \tau$, then $\Gamma \restriction \mathsf{fv}(M) \vdash_{\mathcal{C}} M : \tau$.

Definition 5.7 (Substitution)

We call *type-substitution* to

$$\mathbb{S} = [\tau_1/\alpha_1, \ldots, \tau_n/\alpha_n]$$

where $\alpha_1, \ldots, \alpha_n$ are distinct type variables and $\tau_1, \ldots, \tau_n$ are types in $\mathbb{T}_C$. If τ is a type in $\mathbb{T}_C$, then $\mathbb{S}(\tau)$ is the type obtained by simultaneously substituting α_i by τ_i, with $1 \leq i \leq n$, in τ.

The type $\mathbb{S}(\tau)$ is called an *instance* of the type τ. The notion of substitution can be extended to basis in the following way:

$$\mathbb{S}(\Gamma) \quad = \quad \{x_1 : \mathbb{S}(\tau_1), \ldots, x_n : \mathbb{S}(\tau_n)\} \qquad \text{if } \Gamma = \{x_1 : \tau_1, \ldots, x_n : \tau_n\}$$

The basis $\mathbb{S}(\Gamma)$ is called an *instance* of the basis Γ.

Next we will present some standard properties of the Simple Type System. Details and proofs can be found in [Hindley, 1997].

Lemma 5.8 (Substitution lemmas)

1. If $\Gamma \vdash_{\mathcal{C}} M : \tau$, then $\mathbb{S}(\Gamma) \vdash_{\mathcal{C}} M : \mathbb{S}(\tau)$.
2. If $\Gamma \cup \{x : \tau_1\} \vdash_{\mathcal{C}} M : \tau$ and $\Gamma \vdash_{\mathcal{C}} M : \tau$, then $\Gamma \vdash_{\mathcal{C}} M[N/x] : \tau$.

Theorem 5.9 (Subject reduction)

Let M be a λ-term, and $M \longrightarrow_{\beta}^{*} M'$, then

$$\Gamma \vdash_{\mathcal{C}} M : \tau \Rightarrow \Gamma \vdash_{\mathcal{C}} M' : \tau.$$

The implication in the other direction is called *subject expansion*, and does not hold for this system. For example $(\lambda xy.y)(\lambda z.zz) \longrightarrow_{\beta}^{*} (\lambda y.y)$, where $(\lambda y.y)$ is typable in this system, and $(\lambda xy.y)(\lambda z.zz)$ is not.

Subject reduction also holds for the three sub-calculi defined in Section 2.2.6. The property of subject expansion is verified in the $\lambda_{\mathcal{I}}$ calculus and in the linear and affine λ-calculus.

Theorem 5.10 (Strong normalization)

Let M be a λ-term.

$$\Gamma \vdash_{\mathcal{C}} M : \tau \Rightarrow M \text{ is strongly normalisable.}$$

Notice that the implication in the other direction does not hold. There are many strongly normalisable λ-terms that are not typable in this system. For example, the term $\lambda x.xx$ is in normal form therefore is strongly normalisable, but it is not typable in the Curry Simple Type System (notice that, to type the subterm xx, the variable x has to be of both type α and $\alpha \to \beta$).

5.1.1 Type-checking and Typability

In the Curry Type System, as well in other type systems the following questions arise:

1. Given a term M in Λ, a type τ and a basis Γ, do we have $\Gamma \vdash_{\mathcal{C}} M : \tau$?
2. Given a term M in Λ, is there a type τ and a basis Γ, such that $\Gamma \vdash_{\mathcal{C}} M : \tau$?

The first question concerns *type checking* and the second *typability*, and will be discuss in this section.

Type checking and typability in the Curry Type System are both decidable problems. Moreover, for the typability problem there exists a function that, for

any typable term M, returns the most general type for M in this system. Such type is called the *principal type* of the term.

We first introduce the notions of principal pair and principal type.

Definition 5.11 (Principal pair)

Let M be a term in Λ. Then (Γ, τ) is called a *principal pair* for M if:

1. $\Gamma \vdash_{\mathcal{C}} M : \tau$;
2. If $\Gamma' \vdash_{\mathcal{C}} M : \tau'$, then $\exists\, \mathbb{S}.\ (\mathbb{S}(\Gamma) \subseteq \Gamma'$ and $\mathbb{S}(\tau) \equiv \tau')$.

Note that, if (Γ, τ) is a principal pair for a term M, then $\mathsf{fv}(M) = \mathsf{dom}(\Gamma)$.

Definition 5.12 (Principal type)

Let M be a closed term in Λ. Then τ is called a *principal type* for M if:

1. $\vdash_{\mathcal{C}} M : \tau$;
2. If $\vdash_{\mathcal{C}} M : \tau'$, then $\exists\, \mathbb{S}.\ (\mathbb{S}(\tau) \equiv \tau')$.

The principal type of a term M is a characterisation of the set of types that can be assigned to the term M. Note that every type that can be assigned to M, can be obtained from the principal type of M by applying a type-substitution.

The following result is independently due to Curry [Curry, 1969], Hindley [Hindley, 1969], and Milner [Milner, 1978] (see [Barendregt, 1992] for more details).

Theorem 5.13 (Principal type theorem)

1. Let M be a term in Λ. There exists a total function pp, such that:

   $$\begin{array}{lcl} M \text{ has a type} & \Rightarrow & \mathsf{pp}(M) = (\Gamma, \tau), \text{ and } (\Gamma, \tau) \text{ is a principal pair for } M \\ M \text{ has no type} & \Rightarrow & \mathsf{pp}(M) = \mathsf{fail}. \end{array}$$

2. Let M be a closed term in Λ. There exists a total function pt, such that:

   $$\begin{array}{lcl} M \text{ has a type} & \Rightarrow & \mathsf{pt}(M) = \tau, \text{ and } \tau \text{ is a principal type of } M \\ M \text{ has no type} & \Rightarrow & \mathsf{pt}(M) = \mathsf{fail}. \end{array}$$

Based on the functions define above, decidability for type-checking and typability in the Curry Type System, can be stated. Notice that that for $\Gamma = \{x_1 : \tau_1, \ldots, x_n : \tau_n\}$ we have

$$\Gamma \vdash_{\mathcal{C}} M : \tau \text{ if and only if } \vdash_{\mathcal{C}} \lambda x_1 \ldots x_n.M : \tau_1 \to \cdots \tau_n \to \tau.$$

Therefore one can state the questions of type-checking and type-assignment, taking Γ to be an empty set.

Corollary 5.14

Type checking and typability are decidable problems in the Curry Type System.

Type checking decidability is proved by noticing that, given M and τ:

$$\vdash_{\mathcal{C}} M : \tau \iff \exists\, \mathbb{S}.\ (\mathbb{S}(\tau) = \mathsf{pt}(M)).$$

The type-substitution $\mathbb{S}$ is found using Robinson's of first-order unification [Robinson, 1965], which is decidable.

As for typability, notice that

$$M \text{ is typable} \iff \mathsf{pt}(M) \neq \mathsf{fail}.$$

5.2 Type-inference

We now present two principal type algorithms for the Curry Type System [Wand, 1987, Hindley, 1997], based on Robinson's unification algorithm [Robinson, 1965], that given a term M, return its principal typing.

To type-check a given type for a given term, one can use the type inference algorithms to calculate the principal type of the term and then verify if the given type is an instance of the principal type.

5.2.1 Unification

Robinson's unification algorithm [Robinson, 1965] plays a key role in type inference. We will present a definition of the unification algorithm for the particular case where the terms we want to unify are simple types.

Definition 5.15 (Unifier)

A unifier between two types τ_1 and τ_2, is a type-substitution $\mathbb{S}$ such that $\mathbb{S}\,\tau_1 = \mathbb{S}\,\tau_2$. If two types have a unifier then we say that they are unifiable. We call $\mathbb{S}\,\tau_1$ (or $\mathbb{S}\,\tau_2$), a common instance.

Example 5.16

The types $(\alpha \to \beta \to \alpha)$ and $((\gamma \to \gamma) \to \delta)$ are unifiable. For the substitution $\mathbb{S} = [(\gamma \to \gamma)/\alpha, \beta \to (\gamma \to \gamma)/\delta]$, the common instance is $((\gamma \to \gamma) \to \beta \to (\gamma \to \gamma))$.

Definition 5.17 (Most general unifier)

$\mathbb{S}$ is a most general unifier (mgu) of τ_1 and τ_2 if, for any other unifier $\mathbb{S}_1$, of τ_1 and τ_2, there is a substitutions $\mathbb{S}_2$ such that $\mathbb{S}_1 = \mathbb{S}_2 \circ \mathbb{S}$.

Example 5.18

Consider the types $\tau_1 = (\alpha \to \alpha)$ e $\tau_2 = (\beta \to \gamma)$. The substitution $\mathbb{S}' = [(\alpha_1 \to \alpha_2)/\alpha, (\alpha_1 \to \alpha_2)/\beta, (\alpha_1 \to \alpha_2)/\gamma]$ is a unifier of τ_1 and τ_2, but it is not the mgu.

The mgu of τ_1 and τ_2 is $\mathbb{S} = [\alpha/\beta, \alpha/\gamma]$. The common instance of τ_1 and τ_2 by $\mathbb{S}'$, $(\alpha_1 \to \alpha_2) \to (\alpha_1 \to \alpha_2)$ is an instance of $(\alpha \to \alpha)$.

We now present the Unification algorithm we will use to define type inference. The function UNIFY, given two types τ_1 and τ_2, returns the mgu of τ_1 and τ_2 if it exists, and fails otherwise.

Definition 5.19 (Unification Algorithm)

Let τ_1 and τ_2 be two types. The unification function $\mathsf{UNIFY}(\tau_1, \tau_2)$ is inductively defined as:

$$
\begin{array}{lcl}
\mathsf{UNIFY}(\alpha, \tau) & = & [\tau/\alpha] \ \text{ if } \alpha \notin FV(\tau) \\
 & = & Id \quad \text{(identity function) if } \tau = \alpha \\
 & = & fail \quad \text{otherwise} \\
 & & \\
\mathsf{UNIFY}(\tau_1 \to \tau_2, \alpha) & = & \mathsf{UNIFY}(\alpha, \tau_1 \to \tau_2) \\
 & & \\
\mathsf{UNIFY}(\sigma_1 \to \sigma_2, \tau_1 \to \tau_2) & = & \text{let} \\
 & & \quad \mathbb{S} = \mathsf{UNIFY}(\sigma_2, \tau_2) \\
 & & \text{in } \ \mathsf{UNIFY}(\mathbb{S}\,\sigma_1, \mathbb{S}\,\tau_1) \circ \mathbb{S}
\end{array}
$$

Note that let $\mathbb{S} = \mathsf{UNIFY}(\sigma_2, \tau_2)$ in $\mathsf{UNIFY}(\mathbb{S}\,\sigma_1, \mathbb{S}\,\tau_1) \circ \mathbb{S}$, fails if one of the calls of UNIFY fails.

5.2.2 Milner's Type-Inference Algorithm

The following algorithm, presented in [**?**], defines a function that, given a term M returns a basis Γ and a type τ such that $\Gamma \vdash M : \tau$, is the principal typing of M.

Definition 5.20

Let Γ be a basis, M a term and τ a type. Let UNIFY be the unification function defined above. The function $T(M) = (\Gamma, \tau)$ defines a type inference algorithm for the simply typed λ-calculus, in the following way:

1. If M is a variable x, then $\Gamma = \{x : \alpha\}$ and $\tau = \alpha$, where α is a new variable;
2. If $M \equiv M_1 M_2$, $T(M_1) = (\Gamma_1, \tau_1)$ e $T(M_2) = (\Gamma_2, \tau_2)$, then $T(M) = (\mathbb{S}(\Gamma_1 \cup \Gamma_2), S\alpha)$, where $\mathbb{S} = UNIFY(\tau_1, \tau_2 \to \alpha)$ and α is a new variable;
3. If $M \equiv \lambda x.N$ and $T(N) = (\Gamma_N, \sigma)$ then:
   a) If $x \notin \mathsf{dom}(\Gamma_N)$, then $T(M) = (\Gamma_N, \alpha \to \sigma)$, where α is a new variable;
   b) If $\{x : \tau\} \in \Gamma_N$, $T(M) = (\Gamma_N - \{x : \tau\}, \tau \to \sigma)$.

5.2.3 Wand's Algorithm

The following algorithm [Wand, 1987], given a basis, a term, and a type returns a set of equations, that when applied unification, give the principal type of the term:

Definition 5.21

Let Γ be a basis, $M \in \Lambda$ and $\sigma \in \mathbb{T}$. Let $E = (\Gamma, M, \sigma)$ be the set of equations defined by:

$$\begin{array}{rcl} E(\Gamma, x, \sigma) & = & \{\sigma = \Gamma(x)\} \\ E(\Gamma, MN, \sigma) & = & E(\Gamma, M, \alpha \to \sigma) \cup E(\Gamma, N, \alpha), \\ & & \text{where } \alpha \text{ is a new variable;} \\ E(\Gamma, \lambda x.M, \sigma) & = & E(\Gamma \cup \{x : \alpha\}, M, \beta) \cup \{\alpha \to \beta = \sigma\}, \\ & & \text{where } \alpha, \beta \text{ are new variables.} \end{array}$$

Applying unification to E gives a substitution $\mathbb{S}$ such that $\mathbb{S}(\Gamma) \vdash M : \mathbb{S}(\sigma)$.

5.3 The Damas-Milner Polymorphic Type System

We now present the Damas-Milner [Damas and Milner, 1982] type system for the λ-calculus with parametric polymorphism, which is in the basis of type inference algorithms for functional languages such as ML.

5.3.1 The term language

The set of terms is an extension of the λ-calculus with a a constructor *let*. Given an infinite set of term variables V, the term language is give by the following grammar:

$$M ::= x \texttt{ \textbackslash{} MM'\textbackslash{} } \lambda x.M \texttt{ \textbackslash{} let\textbackslash{} x = M\textbackslash{} in\textbackslash{} M'\$\%\textbackslash{} } fix\, x.M$$

5.3.2 Type Schemes

The definition of types schemes allows us to represent the set of terms one can infer for a term M, given a basis Γ.

Definition 5.22

We say that σ is a type scheme if σ is a simple type τ or a term of the form $\forall_{\alpha_1,...,\alpha_n}.\tau$, where $\alpha_1, ..., \alpha_n$ are called generic type variables.

Definition 5.23

If τ is a type and σ a type scheme, we say that τ is a generic instance of σ if and only if $\sigma = \tau$ or $\sigma = \forall_{\alpha_1,\ldots,\alpha_n}.\eta$ and $\exists_{\tau_1,\ldots,\tau_n}$ such that $\tau = [\tau_i/\alpha_i]\sigma$.

5.3.3 Types

The set of types of this system is given the the following grammar:

$$\begin{array}{lll} \tau & ::= & \alpha \mid \tau' \to \tau'' \\ \sigma & ::= & \tau \mid \forall\alpha.\sigma' \end{array}$$

where α belongs to an infinite set of type variables, τ, τ' and τ'' are simple types, and σ and σ' are type schemes.

5.3.4 The type system

Let x be a term variable, α a type variable, M and M' terms, τ and τ' simple types and σ and σ' type schemes. The Damas-Milner type system is defined by the following inference rules:

$$\text{(axiom)} \qquad \Gamma \cup \{x : \sigma\} \vdash_{ML} x : \sigma$$

$$\text{(gen)} \qquad \frac{\Gamma \vdash_{ML} M : \sigma, \alpha \text{ does not occur free in } \Gamma}{\Gamma \vdash_{ML} M : \forall\alpha.\sigma}$$

$$\text{(inst)} \qquad \frac{\Gamma \vdash_{ML} M : \forall\alpha.\sigma}{\Gamma \vdash_{ML} M : \sigma[\tau/\alpha]}$$

$$\text{(app)} \qquad \frac{\Gamma \vdash_{ML} M : (\tau' \to \tau), \Gamma \vdash_{ML} M' : \tau'}{\Gamma \vdash_{ML} (MM') : \tau}$$

$$\text{(abs)} \qquad \frac{\Gamma \cup \{x : \tau'\} \vdash_{ML} M : \tau}{\Gamma \vdash_{ML} \lambda x.M : (\tau' \to \tau)}$$

$$\text{(let)} \qquad \frac{\Gamma \vdash_{ML} M : \sigma, \Gamma \cup \{x : \sigma\} \vdash_{ML} M' : \sigma'}{\Gamma \vdash_{ML} \text{ let } x = M \text{ in } M' : \sigma'}$$

Example 5.24

Consider the derivation tree for $(\lambda x.x)$ in this system:

$$\{x : \alpha\} \vdash x : \alpha$$
$$\downarrow \text{ABS}$$
$$\vdash \lambda x.x : \alpha \to \alpha$$
$$\downarrow \text{GEN}$$
$$\vdash \lambda x.x : \forall\alpha.\alpha \to \alpha$$

5.3.5 Type Inference

We will define the notion of closure of a type, which will be used in the definition of the algorithm.

Definition 5.25

Let V be a set of type variables, Γ a basis and τ a type. The closure of τ with respect to V, $\overline{V}(\tau)$ is the type scheme $\forall_{\alpha_1,\ldots,\alpha_n}\tau$, where $\alpha_1, \ldots, \alpha_n$ are all the variables that occur in τ that do not belong to V. $\overline{\Gamma}(\tau)$ is the closure of τ with respect to the set of type variables that occur in Γ.

We can now define the type-inference algorithm for this system.

Definition 5.26

Let Γ be a basis, M a term, $\mathbb{S}$ a substitution and τ a type. Let UNIFY be the unification function defined above. The function $W(\Gamma, M) = (\mathbb{S}, \tau)$ is defined in the following way:

1. If M is a variable x and $x : \forall_{\alpha_1,\ldots,\alpha_n}\tau \in \Gamma$ the $\mathbb{S}$ is the identity function and $\sigma = [\beta_i/\alpha_i]\tau$, where β_i is a new variable $(1 \leq i \leq n)$.
2. If $M \equiv M_1 M_2$, let:
   - $W(\Gamma, M_1) = (\mathbb{S}_1, \tau_1)$;
   - $W(\Gamma, M_2) = (\mathbb{S}_2, \tau_2)$;
   - $\mathbb{S}_3 = UNIFY(\mathbb{S}_2\tau_1, \tau_2 \to \beta)$, where β is a new variable;

   then $\mathbb{S} = \mathbb{S}_3 \circ \mathbb{S}_2 \circ \mathbb{S}_1$ and $\sigma = \mathbb{S}\beta$;
3. If $M \equiv \lambda x.N$, let β be a new variable and $W(\Gamma - \{x : \alpha\} \cup \{x : \beta\}, N) = (\mathbb{S}_1, \tau_1)$, then $\mathbb{S} = \mathbb{S}_1$ e $\sigma = \mathbb{S}_1(\beta \to \tau_1)$;

4. If $M \equiv let\ x = M_1\ in\ M_2$, let:
   - $W(\Gamma, M_1) = (\mathbb{S}_1, \tau_1)$ e
   - $W(\mathbb{S}_1(\Gamma - \{x : \alpha\} \cup \{x : \overline{\mathbb{S}_1\Gamma}(\tau_1)\}), M_2) = (\mathbb{S}_2, \tau_2)$

   then $\mathbb{S} = \mathbb{S}_2 \circ \mathbb{S}_1$ e $\sigma = \tau_2$

5.3.6 Correctness and Completeness

Proposition 5.27

If $W(\Gamma, M) = (\mathbb{S}, \tau)$, then $\mathbb{S}(\Gamma) \vdash_{ML} M : \mathbb{S}(\tau)$.

Definition 5.28

Let Γ be a basis and M a term. We say that σ_P is the principal type-scheme of Γ if and only if:

- $\Gamma \vdash_{ML} M : \sigma_P$.
- For any other σ such that $\Gamma \vdash_{ML} M : \sigma$ is a generic instance of σ_P.

Theorem 5.29

Given Γ and M, let Γ' be an instance of Γ and σ a type scheme such that $\Gamma' \vdash_{ML} M : \sigma$. Then:

- $W(\Gamma, M)$ succeeds
- If $W(\Gamma, M) = (\mathbb{S}, \tau)$, then there exists $\mathbb{S}'$ such that $\Gamma' = \mathbb{S}'\mathbb{S}\Gamma$ and σ is a generic instance of $\mathbb{S}'\overline{\mathbb{S}\Gamma}(\tau)$.

6
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